

LTL_f Synthesis Under Environment Specifications

Antonio Di Stasio

Sapienza University of Rome, Italy

VardiFest – FloC 2022



ERC Advanced Grant
WhiteMech:
White-box Self Programming Mechanisms



SAPIENZA
UNIVERSITÀ DI ROMA



SAPIENZA
UNIVERSITÀ DI ROMA

Reactive Synthesis

Given a specification φ over inputs \mathcal{I} and outputs \mathcal{O} , expressed in:

LTL (Pnueli 1977) or LTL_f (De Giacomo, Vardi 2013)

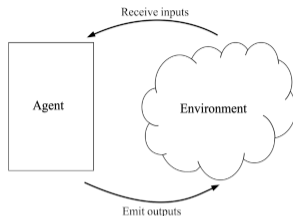
Syntax:

$$\varphi ::= a \mid \varphi \wedge \varphi \mid \neg\varphi \mid \bigcirc\varphi \mid \varphi \mathcal{U} \varphi \mid \diamond\varphi \mid \square\varphi$$

Semantic:

A trace $trace$ is an infinite (LTL) or finite (LTL_f) sequence over \mathcal{I} and \mathcal{O} . We write $trace \models \varphi$ to mean that τ satisfies φ .

Reactive Synthesis



Agent and Environment Strategies, and Traces

For an agent strategy $\sigma_{ag} : \mathcal{I}^+ \rightarrow \mathcal{O}$ and an environment strategy $\sigma_{env} : \mathcal{O}^* \rightarrow \mathcal{I}$, the trace
 $trace(\sigma_{ag}, \sigma_{env}) = (i_1 \cup o_1), (i_2 \cup o_2) \dots \in 2^{\mathcal{I} \cup \mathcal{O}}$
denotes the unique trace induced by both σ_{ag} and σ_{env} .

Synthesis Problem

Given an LTL/ LTL_f task *Goal* for the agent

Find agent strategy σ_{ag} such that $\forall \sigma_{env}. trace(\sigma_{ag}, \sigma_{env}) \models Goal$

Planning (or Synthesis with a model of the world)

Domain

- Planning consider the agent acting in a **(nondeterministic) domain**
- The domain is a **model of how the world** (i.e. the environment) works
- That is, it is a **specification of the possible environment strategies**

$$[[Dom]] = \{\sigma_{env} | \sigma_{env} \text{ compliant with } Dom\}$$

Planning in nondeterministic domains

Given an task *Goal* for the agent, and a domain *Dom* modeling the environment

Find agent behavior σ_{ag} such that $\forall \sigma_{env} \in [[Dom]]. \text{trace}(\sigma_{ag}\sigma_{env}) \models \text{Goal}$

Environment Specifications

Which kinds of environment assumptions can the agent make?

- Nondeterministic planning domains;
- Forms of fairness ($\Box\Diamond\phi$) and stability ($\Diamond\Box\phi$) [ZhuDeGiacomoPuVardiAAAI2020];
- Safety properties [DeGiacomoDiStasioPerelliZhuKR2021];
- GR(1) formulas [DeGiacomoDiStasioTabajaraVardiZhuJCAI2021];
- ..

Environments Specifications as LTL formulas

A natural generalization is to consider general environment specifications expressed as arbitrary LTL formulas [DeGiacomoDiStasioVardiZhuKR2020].

Synthesis Under Environment Specifications

Environment Specifications

Let Env be an LTL/LTL_f formula over $\mathcal{I} \cup \mathcal{O}$.

$$[[Env]] = \{\sigma_{env} \mid \sigma_{env} \text{ satisfies } Env \text{ whatever is the agent strategy}\}$$

Synthesis under environment specifications in LTL/LTL_f

Given an LTL_f task $Goal$ for the agent, and an LTL/LTL_f environment specification Env :

Find agent strategy σ_{ag} such that $\forall \sigma_{env} \in [[Env]]. \text{trace}(\sigma_{ag}, \sigma_{env}) \models Goal$

Theorem [Aminof et al. ICAPS 2019]

To find agent strategy realizing $Goal$ under the environment specification Env , we can use standard synthesis for

$$Env \rightarrow Goal$$

Problem

Solve the synthesis problem for

$$\varphi_{LTL}^e \wedge \varphi_{LTL_f}^e \rightarrow \varphi_{task}^a$$

separating the LTL_f environment specifications:

$$(\varphi_{LTL}^e \wedge \varphi_{LTL_f}^e \rightarrow \varphi_{task}^a) \iff (\varphi_{LTL}^e \rightarrow \varphi_{LTL_f}^e \rightarrow \varphi_{task}^a) \iff (\varphi_{LTL}^e \rightarrow \neg\varphi_{LTL_f}^e \vee \varphi_{task}^a)$$

where $\varphi'_{LTL_f} = \neg\varphi_{LTL_f}^e \vee \varphi_{task}^a$ is expressed in LTL_f and φ_{LTL}^e in LTL.

Key Idea:

- 1 **Agent goal:** $\neg\varphi_{LTL}^e \vee \varphi'_{LTL_f}$ — **Environment goal:** $\varphi_{LTL}^e \wedge \neg\varphi'_{LTL_f}$;
- 2 Build the DFA of φ'_{LTL_f} , solve the reachability game for the **agent** over it;
- 3 Remove the **agent** winning area and do the product with the DPA of φ_{LTL}^e ;
- 4 Solve the parity games for the **environment** over it;
- 5 Combine the two **agent** winning strategies computed in 2 and 4.

Problem

Solve the synthesis problem for

$$\varphi_{GR(1)}^e \rightarrow \varphi_{task}^a$$

Key Idea:

- **Agent goal:** $\neg\varphi_{GR(1)}^e \vee \varphi_{task}^a$ — **Environment goal** $\varphi_{GR(1)}^e \wedge \neg\varphi_{task}^a$
- Build the corresponding DFA $\mathcal{A}_{\varphi_{task}^a}$ of φ_{task}^a , and take its complement $\overline{\mathcal{A}_{\varphi_{task}^a}}$
- Define a GR(1) game whose game arena is $\overline{\mathcal{A}_{\varphi_{task}^a}}$ and winning condition $\varphi_{GR(1)}^e$
- Solve the GR(1) game for the **agent**, i.e., solve the dual of the GR(1) game.

Problem

Solve the synthesis problem for

$$\varphi_S^e \rightarrow \varphi_{task}^a$$

Key Idea:

- Compute the *deterministic safety automaton* \mathcal{D} of φ_S (**no Büchi determinization!**)
- Solve the safety game for the **environment** over \mathcal{D} ;
- Construct the *maximally permissive strategy* \mathcal{T} ;
- Build the corresponding DFA $\mathcal{A}_{\varphi_{task}^a}$ of φ_{task}^a ;
- Do the product of \mathcal{T} and $\mathcal{A}_{\varphi_{task}^a}$;
- Solve the reachability game for the **agent** over it, and return a strategy, if exists.

Problem

Solve the synthesis problem for

$$\varphi_S^e \rightarrow \varphi_{task}^a$$

Key Idea:

- Compute the *deterministic safety automaton* \mathcal{D} of φ_S (**no Büchi determinization!**)
- Solve the safety game for the **environment** over \mathcal{D} ;
- Construct the *maximally permissive strategy* \mathcal{T} ;
- Build the corresponding DFA $\mathcal{A}_{\varphi_{task}^a}$ of φ_{task}^a ;
- Do the product of \mathcal{T} and $\mathcal{A}_{\varphi_{task}^a}$;
- Solve the reachability game for the **agent** over it, and return a strategy, if exists.

No reduction to the implication!

Future Works

Almost all the techniques are based on the following reduction

Theorem [Aminof et al. ICAPS 2019]

To find agent strategy realizing *Goal* under the environment specification *Env*, we can use standard synthesis for

$$Env \rightarrow Goal$$

Possible directions

- In case of safety environment specifications we can directly solve the problem without reduction to the implication.
- What about the other environment specifications? Fairness, $GR(1)$, LTL, ...

Future Works

Almost all the techniques are based on the following reduction

Theorem [Aminof et al. ICAPS 2019]

To find agent strategy realizing *Goal* under the environment specification *Env*, we can use standard synthesis for

$$Env \rightarrow Goal$$

Possible directions

- In case of safety environment specifications we can directly solve the problem without reduction to the implication.
- What about the other environment specifications? Fairness, GR(1), LTL, ...

Moshe's quote: "We need to focus on the problems in *P*, where *P* does not mean Polynomial Time but Practical Problems!"